

Overview of the Cyrix M2 Microprocessor

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Presentation Agenda

- u **Product Goals**
- u **M2 Key Features**
- u **Microarchitecture Overview**
- u **L1 cache and TLB**
- u **Performance**
- u **Multimedia Enhancements**
- u **Summary**

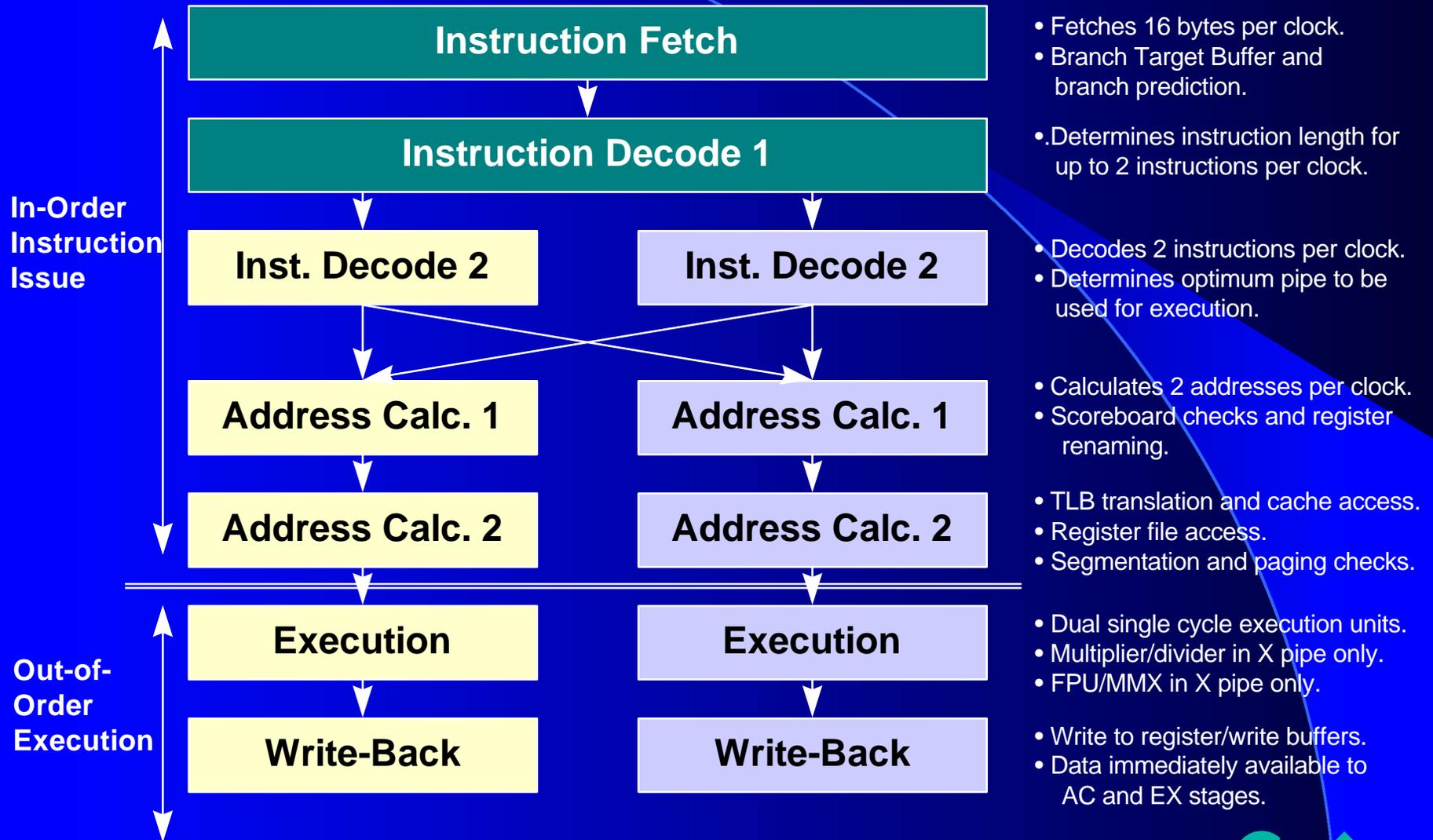
M2 Product Goals

- u Highest performance desktop for Windows[®] 95 and Windows[®] NT in '97**
- u Acceleration of multimedia applications**
- u Provide flexible system cost/performance options**
- u Low cost - highly manufacturable**
- u Low power - mainstream mobile solution**

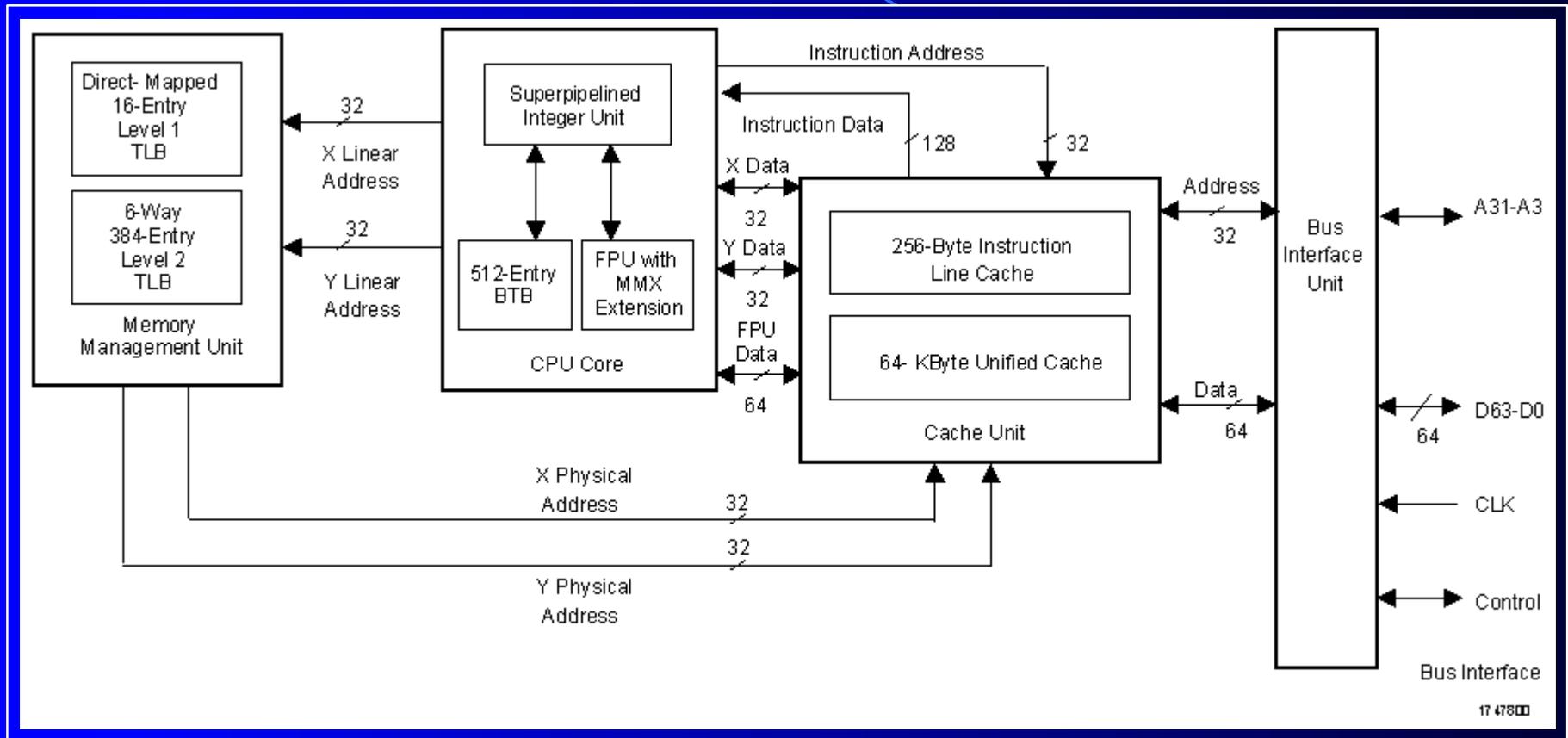
M2 Key Features

- u **High-performance superscalar native x86 core**
 - Optimizations for 16-bit and 32-bit code
 - 2x faster than the 6x86 on 32-bit code
 - Enhanced cache, MMU, BTB
- u **180 to 225 MHz operation**
 - 75/187, 66/200, 75/225 initial offerings
 - 2:1, 2.5:1, 3:1, 3.5:1 core to bus clock ratio
- u **Pentium[®] Pro + MMX instruction set**
- u **Standard 6x86 socket**
 - 2.5 V core, 3.3 V bus interface
 - 6 M transistors
 - less than 200 sq. mm, .35u 5 layer metal CMOS

M2 Pipeline



M2 Block Diagram



M2 Architectural Features

Feature	P6	M2
Full x86 instruction set optimization		√
MMX Instructions		√
Superscalar	√	√
Superpipelined	√	√
Register renaming	√	√
Data dependency removal	√	√
Multi-branch prediction	√	√
Speculative execution	√	√
Out-of-order completion	√	√
80-bit floating point unit	√	√

The M2 is a native mode sixth-generation x86 processor

M2 Branch Unit

- u **Branch target cache**
 - 512 entries
 - 4-way set associative
- u **Branch history table**
 - 512 entries
 - 4-state branch prediction algorithm
- u **Return stack**
 - 8 entries

M2 Memory Management Unit

- u Two level TLB architecture
- u L1 TLB
 - 16 entry, direct mapped
 - Dual ported
- u L2 TLB
 - 384 entry, 6-way set associative
 - Dual ported

M2 Cache Architecture

- u **64 Kbyte unified cache**
 - Primary data cache, secondary instruction cache
 - 4-way set associative, 32 byte line size
 - 2 accesses per cycle
- u **256 byte instruction line cache**
 - Primary instruction cache
 - 8 entries, 32 byte line size
 - Fully associative

M2 Performance

u Windows 95 and Windows NT

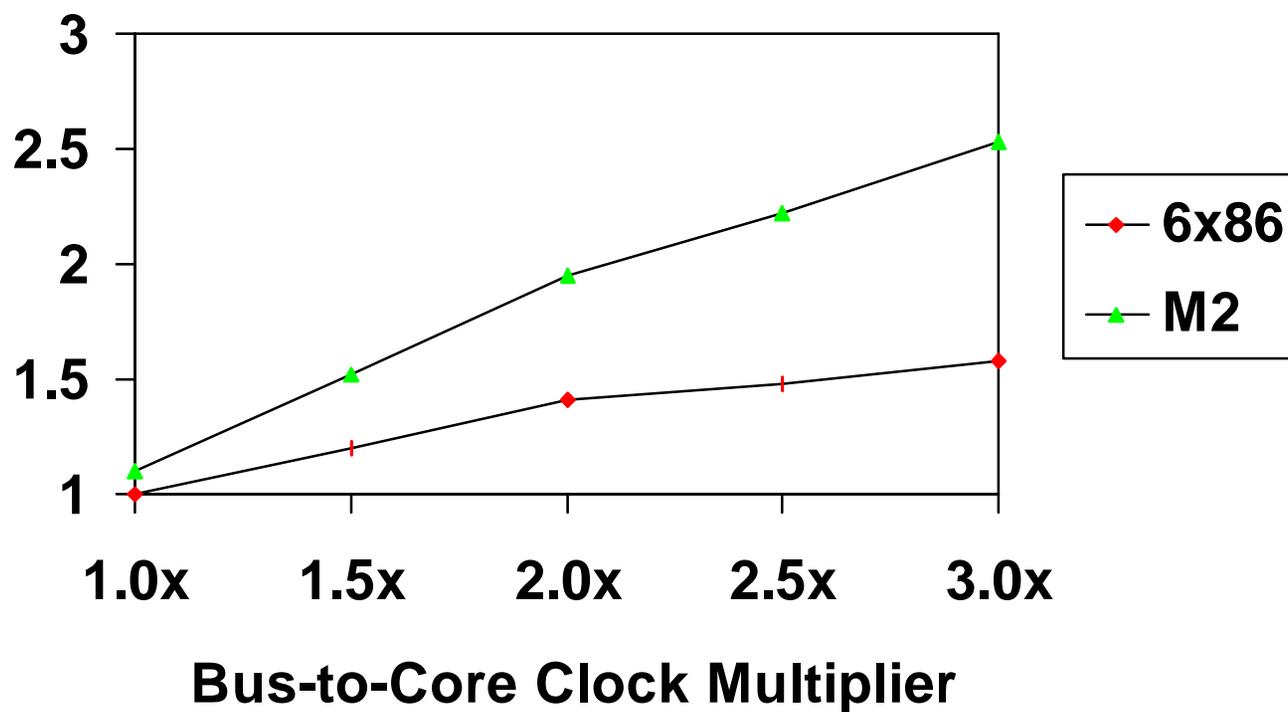
- TLB, cache miss rate increase with 32-bit apps
- Critical word latency limits performance

	Pentium Pro	M2
L1 Cache	8K instr. 8K data	64K unified
L2 Cache	On module 256K or 512K	Off chip 256K or 512K
TLB	32-entry instr. 64-entry data	16-entry L1 384-entry L2
BTB	512-entry	512-entry

- ## u Growing mismatches between processor frequency and bus frequency exacerbate the cache miss penalty
- 1 bus cycle = 2-3 processor cycles

M2 Performance Analysis

Relative Performance - 32-bit Applications



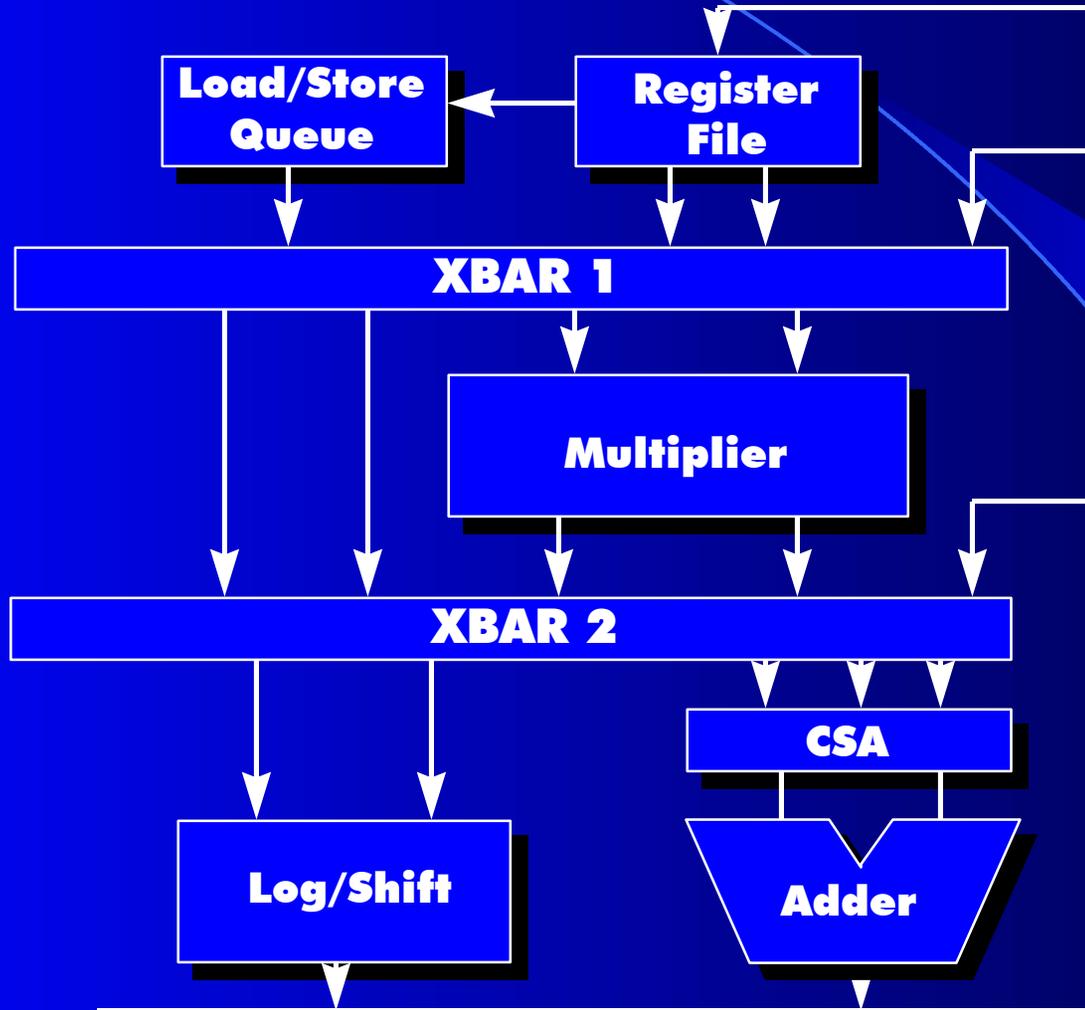
M2 Multimedia Extensions: Overview

- u **New data types: packed Byte/Word/Doubleword, Quadword**
- u **Eight media registers**
 - Aliased with floating point (FP) register state
 - 64-bit general purpose registers
- u **Instruction set**
 - Single Instruction Multiple Data (SIMD)
 - Saturation or modulo arithmetic
- u **Instructions**
 - Arithmetic, Comparison, Conversion, Logical, Shift, & Data Transfer
 - Source operand can reside in memory or in media register
 - Destination operand can only reside in media register

Integrated Multimedia & FPU

- u Existing wide adders & multiplier in FP module can be efficiently subdivided to support SIMD processing
- u FPU pipeline design can accommodate single-cycle multiply & fused multiply-add operations
- u MMX compatibility prohibits simultaneous use of FPU hardware by both x86 FP & MMX operations
 - Why add dedicated functional unit/hardware?

M2 Multimedia/FPU Block Diagram



M2 Multimedia Architecture: Pipeline Diagram

Multimedia instructions execute in a variable length pipeline. Instructions sent to instruction shelf in IQ stage.

Arithmetic/Logical/Shift:



IF ==> Fetch
ID1 ==> Instruction Decode 1
ID2 ==> Instruction Decode 2
AC1 ==> Address Calculation
AC2 ==> Address Calculation/Cache Access
IQ ==> Transfer to MMX Instruction Shelf
RF ==> FP Register File Access
ALS ==> Arithmetic/Logical/Shift
M1 ==> Multiply Stage 1
M2 ==> Multiply Stage 2
WB ==> Write-Back

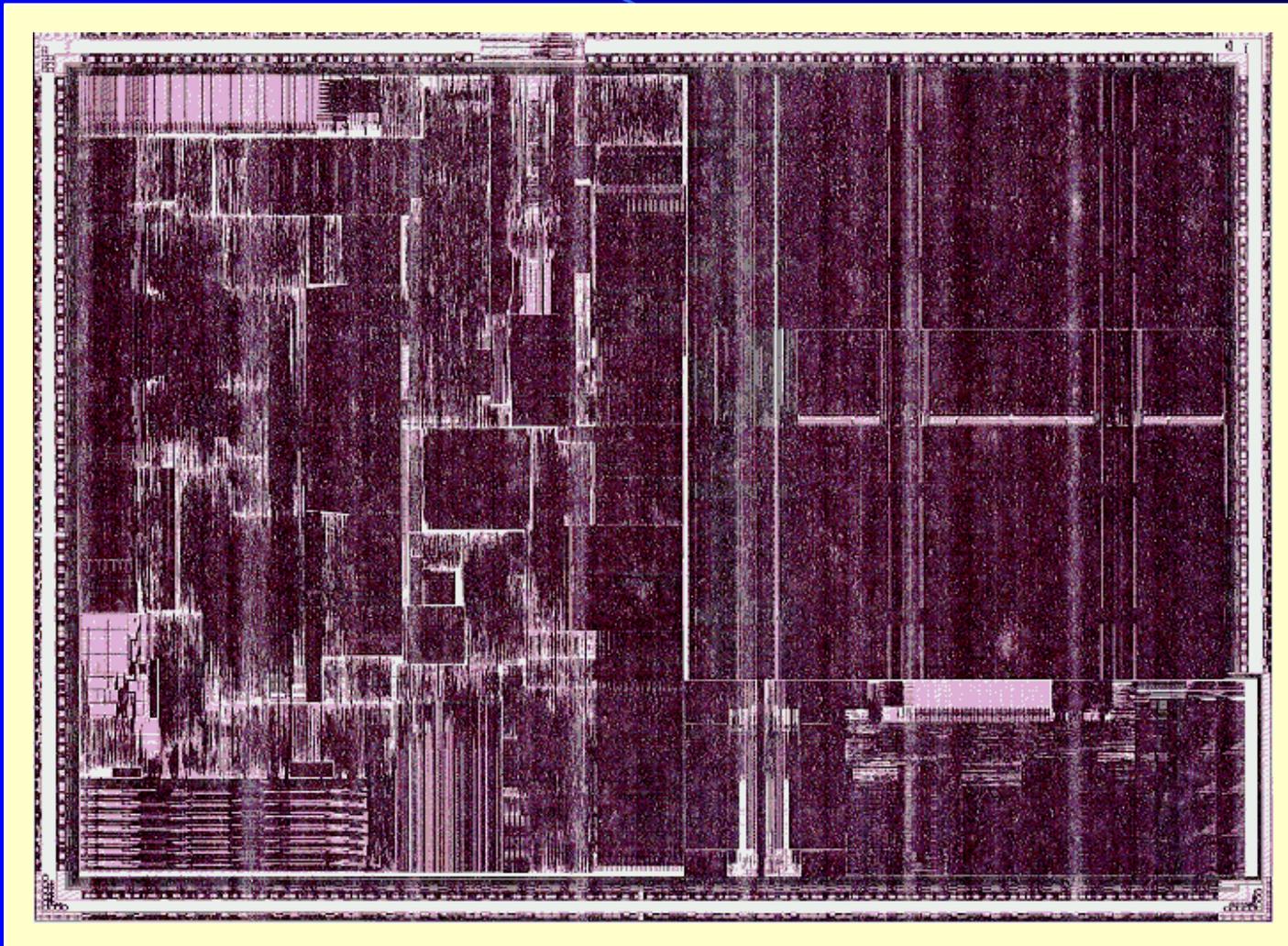
Multiply Pipeline:



M2 Multimedia Performance

- u **Cache line locking ==> scratch pad memory**
 - Locked memory lines guarantee locality of reference
 - Used by driver software code & data
 - Predictable access speed yields real-time capability
- u **Pipeline accesses to L1 cache**
 - Memory operand access at same speed as register access
 - Combined with lockdown capability, permits giant “register files”
- u **Single-cycle execution**
 - EMMS, Add, Subtract, Logical, & Shift operations execute in a single cycle
 - Multiply & MAC execute with single-cycle throughput & latency of 2 clocks

M2 Die



Summary

- u High-performance superscalar x86 microprocessor
- u Optimizations for 16-bit and 32-bit code
- u 200 MHz core frequency
- u MMX software compatible
- u Utilizes existing board and chipset infrastructure
- u Production: 1H '97